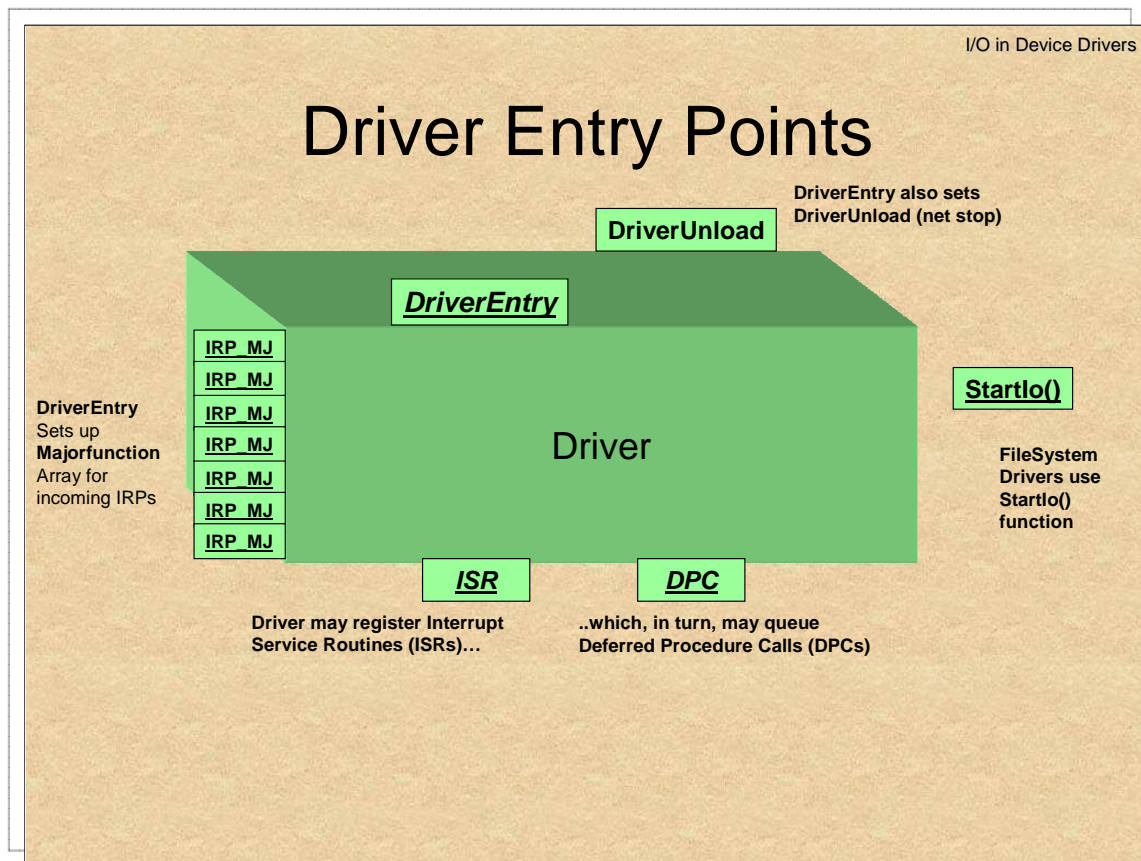


I/O in Device Drivers

Creating a Driver that actually DOES something

This section extends the sample driver by introducing I/O processing via IRPs, and interacting with the driver from user mode and from other drivers.

Key Concepts: IRP, IRP Dispatching, Buffered I/O, Direct I/O, IoControlCodes (IOCTLs)



The Kernel defines two callback interfaces for drivers:

Fast I/O

Rapid synchronous I/O only, mostly for File System Drivers
Direct from user buffers to system cache (less copying)

I/O Request Packets

Default I/O for most operations:

Both synchronous and asynchronous I/O
Page faults implemented by IRPs to file system
Networking – send/recv implemented as IRPs

Driver may define additional entry points/callbacks. Fast I/O is used primarily for File System Drivers (FSDs), and is left out of the scope of this course.

IRPs

- I/O operations are put into “I/O Request Packets”
- IRPs pass up and down the driver stack
- Every driver owns an “IO_STACK_LOCATION” in IRP
- Top level (creator of IRP) must set up IRP “stack size”
- Structure documented, but remains semi-opaque
 - Structs of Unions of Structs – **very** volatile

A fundamental concept in the Windows I/O architecture is that of an **I/O Request Packet**, or **IRP**.

IRPs - I/O Request Packets

- IRP_MJ: “Major” Requests
- IRP_MN: “Minor” (sub) Requests (e.g. for IRP_MJ_PNP)
- Common Major request types:

IRP_MJ_	Use
CREATE	File/Socket/Dir creation open
CLOSE	File/Socket/Dir close
DEVICE_CONTROL	Ioctl/DeviceIoControl
FILESYSTEM_CONTROL	Various FSD operations
READ	Read operation
QUERY_INFORMATION	Get information on descriptor
SET_INFORMATION	Set information of descriptor
WRITE	Write operation

Kernel drivers (with the exception of NDIS and FSD) generally communicate through I/O Request Packets. These “packets” are semi opaque objects.

The Kernel defines IRP_MJ_ types, corresponding to “Major” codes, and IRP_MN_ types, corresponding to “Minor” codes.

The Major codes are for the various request operations, the important ones of which are shown above. The Minor codes are sub codes for a particular Major – for example, Plug and Play operations all have the same Major code, IRP_MJ_PNP, but specific minor codes for starting/stopping devices, etc.

```
typedef struct _IRP {
    ...
    PMDL MdlAddress;
    ULONG Flags;
    union {
        ...
        PVOID SystemBuffer;
    } AssociatedIrp;
    ...
    IO_STATUS_BLOCK IoStatus;
    KPROCESSOR_MODE RequestorMode;
    ...
    BOOLEAN Cancel; // The cancel bit
    ...
    PDRIVER_CANCEL CancelRoutine;
    PVOID UserBuffer;
    union {
        struct { ..
            union {
                KDEVICE_QUEUE_ENTRY DeviceQueueEntry;
                struct {
                    PVOID DriverContext[4];
                };
            };
        };
    };
    ...
    PETHREAD Thread;
    LIST_ENTRY ListEntry;
    .. } Overlay;
    } Tail;
} IRP, *PIRP;
```

Type	Size
MdlAddress	
Flags	
AssociatedIrp	
ThreadListEntry	
IoStatusBlock	
RequestorMode	PendingReturned
StackCount	CurrentLocation
Cancel	CancelIrql
ApcEnvironment	AllocationFlags
UserIoSb	
UserEvent	
Overlay	
CancelRoutine	
UserBuffer	
Tail	

IRPs

Type: Specifies this structure to be an IRP. Reserved.

Size: sizeof (struct IRP) +
StackCount * sizeof(IO_STACKLOCATION)

Flags: Read-only for File System Drivers

IRP_NOCACHE, IRP_PAGING_IO, IRP_MOUNT_COMPLETION
IRP_SYNCHRONOUS_API, IRP_ASSOCIATED_IRP,
IRP_BUFFERED_IO, IRP_DEALLOCATE_BUFFER
IRP_INPUT_OPERATION, IRP_SYNCHRONOUS_PAGING_IO
IRP_CREATE_OPERATION, IRP_READ_OPERATION
IRP_WRITE_OPERATION, IRP_CLOSE_OPERATION
IRP_DEFER_IO_COMPLETION

RequestorMode: KernelMode or UserMode

PendingReturned: IoMarkIrpPending

Cancel: IRP has been canceled

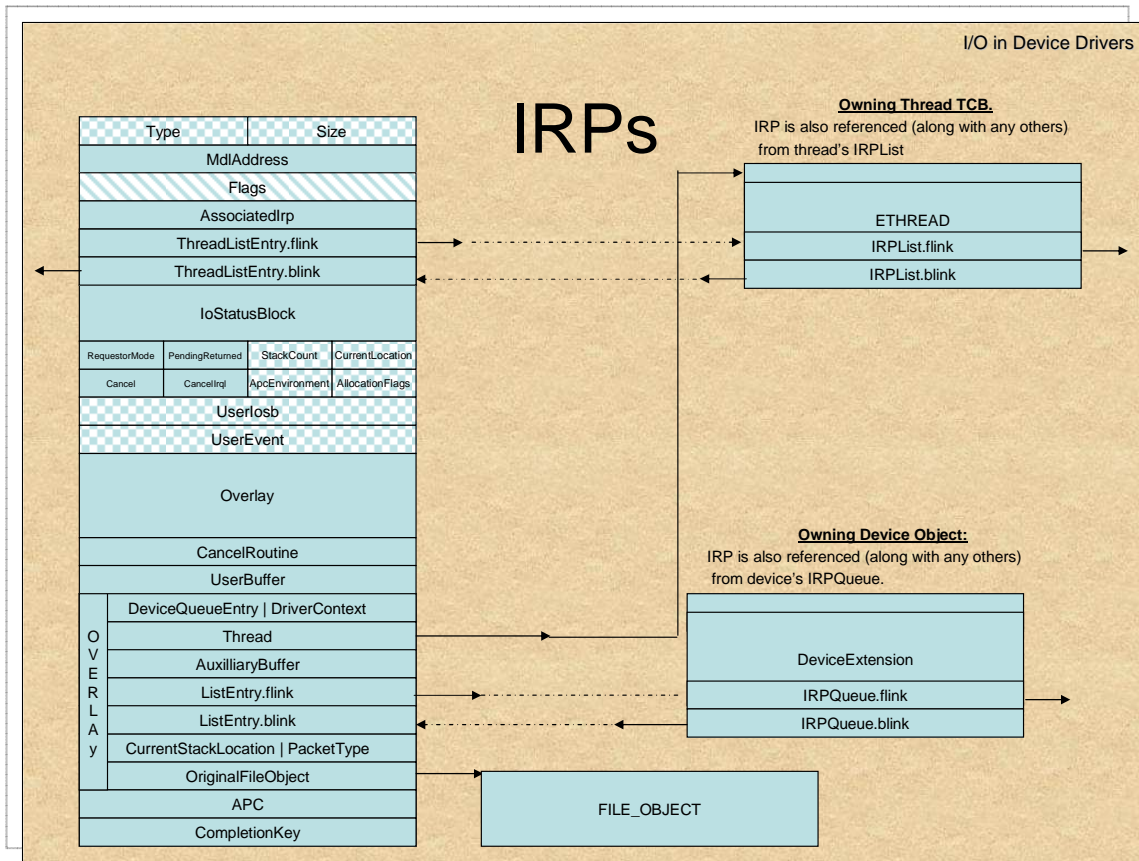
CancelRoutine: IRP Cancel Routine set by driver

Overlay: APC associated with this IRP

Tail: Structures for Kernel Use

The illustrations on the next three pages show the IRP structure. Notice how it is aligned in memory for fast access.

It should be noted that most of these fields are NOT to be handled directly, and there exist functions and macros for that.



Lastly, the IRP “Tail” is a union containing all the fields that the IRP cannot hope to align. Most of these fields are in an “overlay” struct, and they link the IRP to its corresponding device and thread.

```

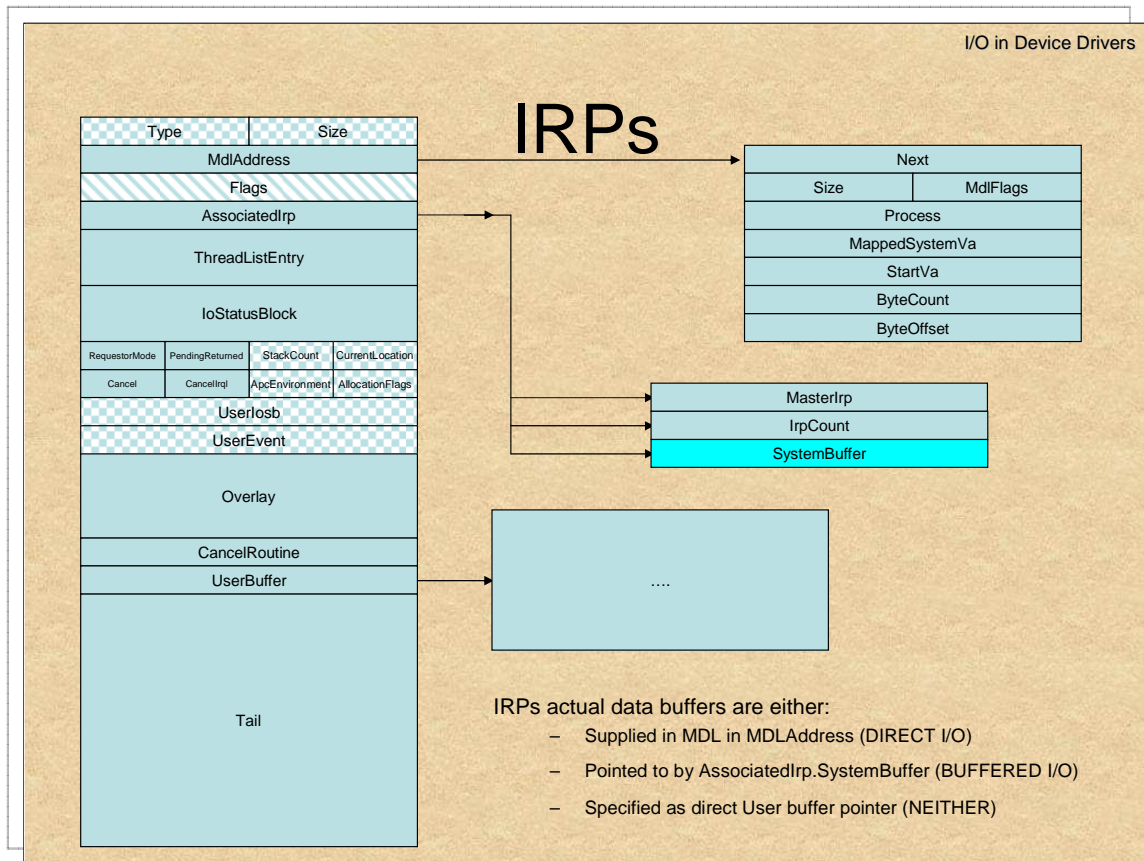
union {
    struct {
        union {
            KDEVICE_QUEUE_ENTRY DeviceQueueEntry;
            struct {
                PVOID DriverContext[4]; } ;
        } ;

        PETHREAD Thread;
        PCHAR AuxiliaryBuffer;

        struct {
            LIST_ENTRY ListEntry;
            union {
                struct _IO_STACK_LOCATION *CurrentStackLocation;
                ULONG PacketType;
            };
        };
        PFILE_OBJECT OriginalFileObject;
    } Overlay;

    KAPC Apc;
    PVOID CompletionKey;
} Tail;

```



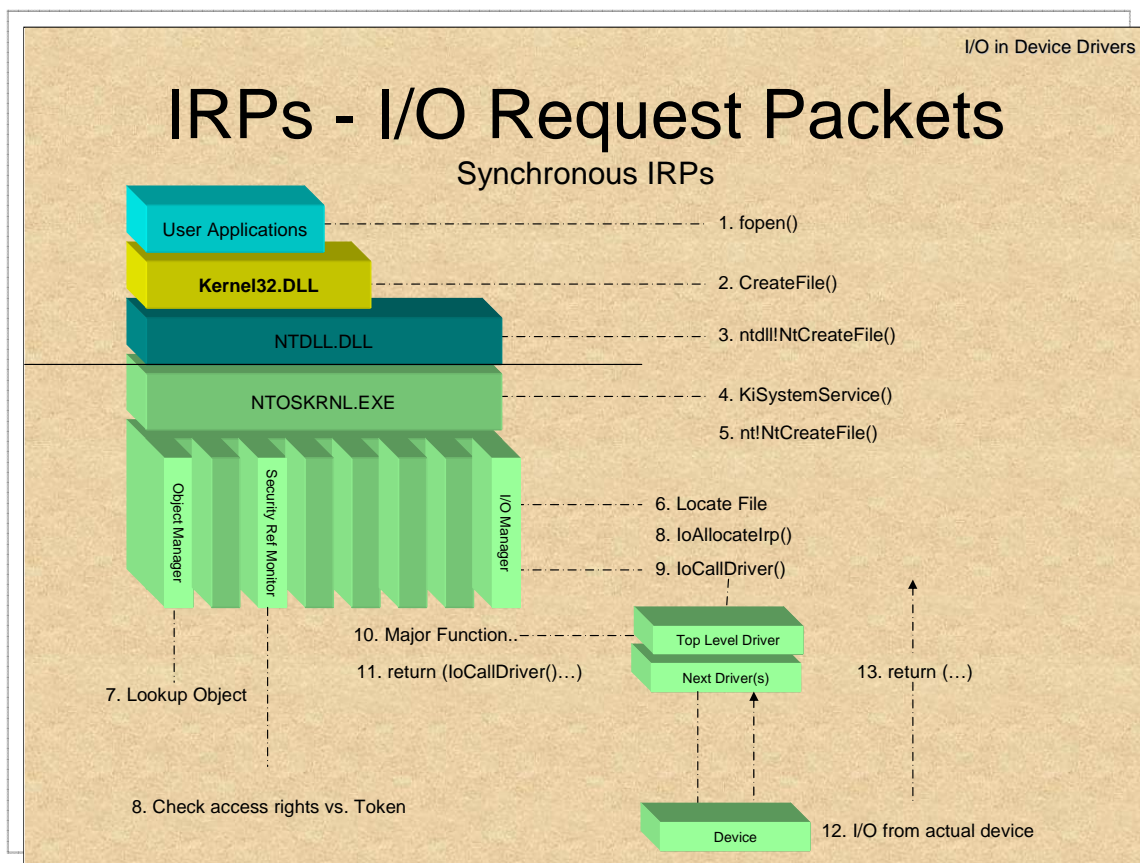
IRPs relate to I/O requests, and therefore point to data buffers. It's not that simple, however, as there are three modes of access to data buffers:

DIRECT I/O: In which the IRP contains a pointer to an MDL (in the MdlAddress field, as shown above). This MDL contains the virtual pages associated with the IRP, and it is the device driver's responsibility to lock these pages in memory.

BUFFERED I/O: In which the IRP contains a pointer to locked in memory pages – the I/O manager takes care of all the lock operations, etc. However, this involves buffering and therefore an extra copy operation. So, while it is easier to handle, it is also more expensive performance-wise.

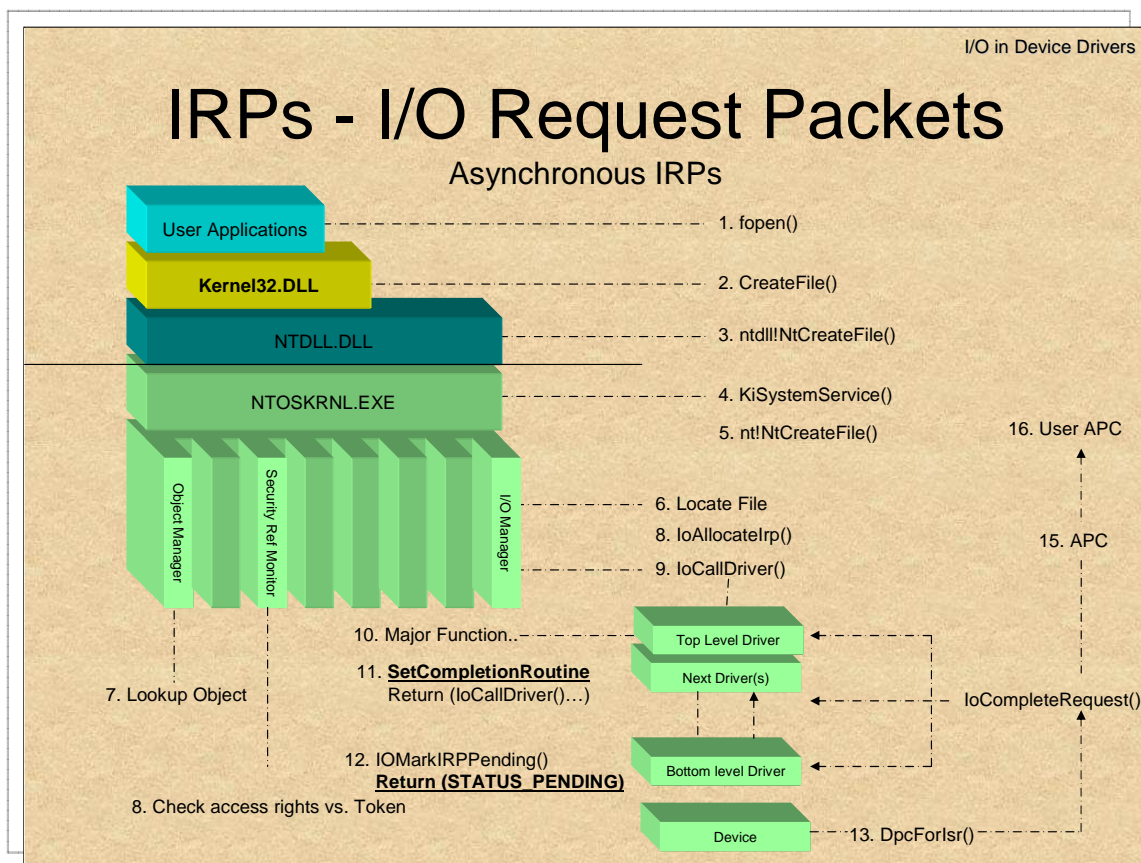
NEITHER: Used in I/O Control codes (IOCTLS, described later), this mode simply passes the user buffer address to the driver. The driver needs to prepare and handle the MDL.

The Device object specifies the preferred mode of operation in its Flags (see page 80). IOCTLS, however, may use any of the modes.



After getting familiar with the IRP structure, we can next look at the typical flow of an I/O request – from inception (usually, in user mode) down to the device.

This slide depicts the process of such a typical I/O request – in this case a file “fopen()” from the C standard library. This is a synchronous request, meaning the process blocks until the I/O returns.



Processing of Asynchronous IRPs is similar all the way up to the first driver called by the I/O manager. Drivers determining a request to be potentially asynchronous still push the IRP down the stack, but may each opt to set an IRP Completion Routine of their own.

When the bottom driver processes the request, it usually submits a request to the hardware device, `IoMarkIrpPending()` and returns a `STATUS_PENDING`. This bubbles up the driver chain back to the I/O manager.

The IRP is completed in a truly asynchronous manner when some other entity chooses to call **`IoCompleteRequest()`** on it. This function is usually called from the device driver who set the interrupt handler (ISR) on the device. **`IoCompleteRequest()`** then calls the completion routines of the IRP, in reverse order. Finally, it signals the I/O manager that the IRP is indeed complete, which in turn schedules an Asynchronous Procedure Call (APC), and completes any user APCs that may have been scheduled, as well.

Handling IRPs

- Drivers implement IRP dispatcher callback functions

```
pDriverObject->MajorFunction[IRP_DISPATCHED] = ... ;
```

- All dispatchers implement the same interface

```
NTSTATUS DispatcherName(IN PDEVICE_OBJECT DeviceObject,
                      IN PIRP Irp);
```

- Simplest implementation: One dispatcher, switch()

The DRIVER_OBJECT struct has an array of **MajorFunctions**, with indices corresponding to all the major IRP types shown so far – that is, the IRP_MJ constants, which are actually implemented as an enum.

To handle IRPs, we would have our driver main function look something like this:

```
// Prototype - We'll implement this later..
NTSTATUS driverIRPDispatcher(PDEVICE_OBJECT DeviceObject, PIRP Irp);

// Our DriverEntry function:
NTSTATUS DriverEntry (IN PDRIVER_OBJECT pDriverObject,
                    IN PUNICODE_STRING strRegistryPath )
{
    pDriverObject->DriverUnload = driverCleanupFunction;

    // Register Callback some IRP, say, IRP_MJ_READ
    pDriverObject->MajorFunction[IRP_MJ_READ] = driverIRPDispatcher;

    DbgPrint("Driver:: Hello, Kernel!\n");
    return STATUS_SUCCESS;
}
```

Listing 1: Stub Driver, Entry Function, IRP Dispatcher registration

Since all dispatchers implement the same prototype, it's often simpler to just implement one, and register it to multiple IRPs, perhaps even all of them. To register more than one IRP to the same IRP Dispatcher you could use something like:

```
pDriverObject->MajorFunction[IRP_TO_PROCESS] =
    ..
pDriverObject->MajorFunction[IRP_TO_ALSO_PROCESS] = driverIRPDispatcher;
```

Not all IRPs must be handled. As the following example shows, IRPs that do not have installed handlers are handled by `IoPInvalidDeviceRequest`, which returns an error code for the IRP.

```
kd> !drvobj kmixer 7
Driver object (8697a988) is for:
  \Driver\kmixer
Driver Extension List: (id , addr)

Device Object list:
Device object (8653f528) is for:
  \Driver\kmixer DriverObject 8697a988
Current Irp 00000000 RefCount 0 Type 0000002f Flags 00002010
DevExt 8653f5e0 DevObjExt 8653f5e8
ExtensionFlags (0000000000)
AttachedTo (Lower) 866d9570 \Driver\swenum
Device queue is not busy.

DriverEntry:   b7203105   kmixer!GsDriverEntry
DriverStartIo: 00000000
DriverUnload:  b71ea610   kmixer!DriverUnload

Dispatch routines:
[00] IRP_MJ_CREATE                f7039fe2   ks!DispatchCreate
[01] IRP_MJ_CREATE_NAMED_PIPE    804fb8de   nt!IopInvalidDeviceRequest
[02] IRP_MJ_CLOSE                 f7039711   ks!DispatchClose
[03] IRP_MJ_READ                  804fb8de   nt!IopInvalidDeviceRequest
[04] IRP_MJ_WRITE                 f70391cc   ks!DispatchWrite
[05] IRP_MJ_QUERY_INFORMATION     804fb8de   nt!IopInvalidDeviceRequest
[06] IRP_MJ_SET_INFORMATION       804fb8de   nt!IopInvalidDeviceRequest
[07] IRP_MJ_QUERY_EA              804fb8de   nt!IopInvalidDeviceRequest
[08] IRP_MJ_SET_EA                804fb8de   nt!IopInvalidDeviceRequest
[09] IRP_MJ_FLUSH_BUFFERS        804fb8de   nt!IopInvalidDeviceRequest
[0a] IRP_MJ_QUERY_VOLUME_INFORMATION 804fb8de   nt!IopInvalidDeviceRequest
[0b] IRP_MJ_SET_VOLUME_INFORMATION 804fb8de   nt!IopInvalidDeviceRequest
[0c] IRP_MJ_DIRECTORY_CONTROL    804fb8de   nt!IopInvalidDeviceRequest
[0d] IRP_MJ_FILE_SYSTEM_CONTROL   804fb8de   nt!IopInvalidDeviceRequest
[0e] IRP_MJ_DEVICE_CONTROL        f7038f60   ks!DispatchDeviceIoControl
[0f] IRP_MJ_INTERNAL_DEVICE_CONTROL 804fb8de   nt!IopInvalidDeviceRequest
[10] IRP_MJ_SHUTDOWN              804fb8de   nt!IopInvalidDeviceRequest
[11] IRP_MJ_LOCK_CONTROL          804fb8de   nt!IopInvalidDeviceRequest
[12] IRP_MJ_CLEANUP               804fb8de   nt!IopInvalidDeviceRequest
[13] IRP_MJ_CREATE_MAILSLLOT     804fb8de   nt!IopInvalidDeviceRequest
[14] IRP_MJ_QUERY_SECURITY        804fb8de   nt!IopInvalidDeviceRequest
[15] IRP_MJ_SET_SECURITY          804fb8de   nt!IopInvalidDeviceRequest
[16] IRP_MJ_POWER                 f70327cf   ks!KsDefaultDispatchPower
[17] IRP_MJ_SYSTEM_CONTROL        b72014d0   kmixer!PerfWmiDispatch
[18] IRP_MJ_DEVICE_CHANGE         804fb8de   nt!IopInvalidDeviceRequest
[19] IRP_MJ_QUERY_QUOTA           804fb8de   nt!IopInvalidDeviceRequest
[1a] IRP_MJ_SET_QUOTA             804fb8de   nt!IopInvalidDeviceRequest
[1b] IRP_MJ_PNP                   b71ea570   kmixer!DispatchPnp
```

IRP Debugging

Useful Tools:

KD:

!irpfind: Find all active IRPs in the system

!irp: Display IRP specific data

OSR: IRPTracker

The Kernel Debugger offers powerful extensions for diagnosing and debugging IRPs. The first is “!irpfind”, that searches the non-paged pool for memory allocations with a tag of “Irp”, and then walks them, and provides summary data:

```
kd> !irpfind
Searching NonPaged pool (81337000 : 82400000) for Tag: Irp?

   Irp      [ Thread ]  irpStack: (Mj,Mn)   DevObj  [Driver]           MDL Process
81d674e8 [81f24558] irpStack: ( e,2d)  821612a0 [ \Driver\AFD]
81e2eb28 [81f23368] irpStack: ( e,2d)  821612a0 [ \Driver\AFD]
81e332b0 [8226dda8] irpStack: ( c, 2)  8232e020 [ \FileSystem\Ntfs]
81e3c008 [00000000] Irp is complete (CurrentLocation 3 > StackCount 2) 0x823b4788
81e3c528 [00000000] Irp is complete (CurrentLocation 3 > StackCount 2) 0x823b4788
81e3cd78 [00000000] Irp is complete (CurrentLocation 3 > StackCount 2) 0x823b4788
81e63570 [81e3d560] irpStack: ( c, 2)  8232e020 [ \FileSystem\Ntfs]
81e6a238 [82336a08] irpStack: ( e,43)  821612a0 [ \Driver\AFD]
81e70528 [00000000] Irp is complete (CurrentLocation 3 > StackCount 2) 0x823b4788
81e70950 [00000000] Irp is complete (CurrentLocation 3 > StackCount 2) 0x823b4788
81e70be0 [00000000] Irp is complete (CurrentLocation 3 > StackCount 2) 0x823b4788

81e8c3a0 [00000000] irpStack: ( 0, 0)  823078a8 [ \Driver\Cdrom]
```

The table lists the IRPs found, their owning threads (a PETHREAD), the owning device object, Device Driver, IRP Major and Minor code, and MDL, if any.

Specific detail for a particular IRP can then be displayed using “!irp” on the IRP address:

```
kd> !irp 81e332b0
Irp is active with 8 stacks 8 is current (= 0x81e3341c)
No Mdl: No System Buffer: Thread 8226dda8: Irp stack trace.
  cmd  flg  cl  Device      File      Completion-Context
[  0,  0]  0  0  00000000  00000000  00000000-00000000

                                     Args: 00000000 00000000 00000000 00000000
.....
>[  c,  2]  1  1  8232e020  822f3400  00000000-00000000      pending
                                     \FileSystem\Ntfs
                                     Args: 00000020 00000017 00000000 00000000
```

As well as, of course, other commands, like !thread and !devobj. !thread is especially useful, as it shows the thread's entire IRPList, as well as the process name.

```
kd> !thread 8226dda8
THREAD 8226dda8  Cid 0670.06ac  Teb: 7ffd6000 Win32Thread: e195aad0 WAIT:
.....
IRP List:
 82297248: (0006,0190) Flags: 00000000  Mdl: 00000000
 82261b68: (0006,0190) Flags: 00000000  Mdl: 00000000
 81ea9dd8: (0006,0190) Flags: 00000000  Mdl: 00000000
 82035e70: (0006,0190) Flags: 00000000  Mdl: 00000000
 81e41008: (0006,0190) Flags: 00000000  Mdl: 00000000
 822db3e0: (0006,0190) Flags: 00000000  Mdl: 00000000
 81e332b0: (0006,0190) Flags: 00000000  Mdl: 00000000

Not impersonating
DeviceMap                e1cc4470
Owning Process            0          Image:          <Unknown>
Attached Process       820cbda0      Image:        explorer.exe
.....
```

DT'ing also helps:

```
kd> dt !_IRP 81e332b0
ntdll!_IRP
+0x000 Type                : 6
+0x002 Size                 : 0x190
+0x004 MdlAddress           : (null)
+0x008 Flags                 : 0
+0x00c AssociatedIrp         : __unnamed
+0x010 ThreadListEntry      : _LIST_ENTRY [ 0x8226dfb8 - 0x822db3f0 ]
+0x018 IoStatus              : _IO_STATUS_BLOCK
+0x020 RequestorMode         : 1 ''
+0x021 PendingReturned      : 0 ''
+0x022 StackCount           : 8 ''
+0x023 CurrentLocation      : 8 ''
+0x024 Cancel                : 0 ''
+0x025 CancelIrql           : 0 ''
+0x026 ApcEnvironment       : 0 ''
+0x027 AllocationFlags      : 0xc ''
+0x028 UserIosb              : 0x7c8837e0 _IO_STATUS_BLOCK
+0x02c UserEvent             : (null)
+0x030 Overlay               : __unnamed
+0x038 CancelRoutine         : 0x80512601      void nt!FsRtlCancelNotify+0
+0x03c UserBuffer            : 0x7c883800
+0x040 Tail                  : __unnamed
```

Because, remembering that “Tail” contains many useful parameters the Kernel associates this IRP with, one can quickly deduce:

```
kd> dd 81e332b0 + 0x40
81e332f0  00000000 00000000 00000000 00000000
81e33300  8226dda8 00000000 e187ab68 e187ab68
81e33310  81e3341c 822f3400 00000000 00000000
81e33320  00000000 00000000 00000000 00000000
81e33330  00000000 00000000 00000000 00000000
81e33340  00000000 00000000 00000000 00000000
81e33350  00000000 00000000 00000000 00000000
81e33360  00000000 00000000 00000000 00000000
lkd> dt _FILE_OBJECT 822f3400
ntdll!_FILE_OBJECT
+0x000 Type           : 5
+0x002 Size           : 112
.....
+0x02c Flags          : 0x40000
+0x030 FileName       : _UNICODE_STRING "\Docume~1\All Users\Desktop"
.....
```

For real time statistics, either attach a Kernel Debugger, or use OSR’s “IRPTracker” Utility. The figure below shows a capture of a “type C:\temp.txt” command from cmd.exe.

The screenshot shows the OSR's IrpTracker Utility V2.18 window. The main window displays a table of IRP operations. The table has columns for Time, Call/Comp, IRP Addr-Seq Number, Originating Device, Target Device, Major Function, Minor Function, and Completion Status. The operations are listed in chronological order, showing various file system operations performed by cmd.exe.

Time	Call/Comp	IRP Addr-Seq Number	Originating Device	Target Device	Major Function	Minor Function	Completion Status
04:04:50.261	NTAPI	NtOpenFile	cmd.exe	(0x8207E950) sr	CREATE		
04:04:50.261	Call	0x81F916D8-4		(0x8232E020) Ntfs	CREATE		
04:04:50.261	Comp	0x81F916D8-4		(0x8232E020) Ntfs	CREATE		SUCCESS, Inf
04:04:50.261	NTAPIRet	NtOpenFile	cmd.exe	(0x8207E950) sr	CREATE		SUCCESS, Inf
04:04:50.261	NTAPI	NtQueryDirectoryFile	cmd.exe	(0x8207E950) sr	DIRECTORY_CONTROL	QUERY_DIRE...	
04:04:50.261	Call	0x81F916D8-5		(0x8232E020) Ntfs	DIRECTORY_CONTROL	QUERY_DIRE...	
04:04:50.261	Comp	0x81F916D8-5		(0x8232E020) Ntfs	DIRECTORY_CONTROL	QUERY_DIRE...	SUCCESS, Inf
04:04:50.261	NTAPIRet	NtQueryDirectoryFile	cmd.exe	(0x8207E950) sr	DIRECTORY_CONTROL	QUERY_DIRE...	SUCCESS, Inf
04:04:50.261	NTAPI	NtCreateFile	cmd.exe	(0x8207E950) sr	CREATE		
04:04:50.261	Call	0x81F916D8-6		(0x8232E020) Ntfs	CREATE		
04:04:50.261	Comp	0x81F916D8-6		(0x8232E020) Ntfs	CREATE		SUCCESS, Inf
04:04:50.261	NTAPIRet	NtCreateFile	cmd.exe	(0x8207E950) sr	CREATE		SUCCESS, Inf
04:04:50.261	NTAPI	NtQueryVolumeInform...	cmd.exe	(0x8207E950) sr	QUERY_VOLUME_INFO...		
04:04:50.261	NTAPIRet	NtQueryVolumeInform...	cmd.exe	(0x8207E950) sr	QUERY_VOLUME_INFO...		SUCCESS, Inf
04:04:50.261	NTAPI	NtQueryVolumeInform...	cmd.exe	(0x8207E950) sr	QUERY_VOLUME_INFO...		
04:04:50.261	NTAPIRet	NtQueryVolumeInform...	cmd.exe	(0x8207E950) sr	QUERY_VOLUME_INFO...		SUCCESS, Inf
04:04:50.261	NTAPI	NtQueryInformationFile	cmd.exe	(0x8207E950) sr	QUERY_INFORMATION		
04:04:50.261	NTAPIRet	NtQueryInformationFile	cmd.exe	(0x8207E950) sr	QUERY_INFORMATION		SUCCESS, Inf
04:04:50.261	NTAPI	NtSetInformationFile	cmd.exe	(0x8207E950) sr	SET_INFORMATION		
04:04:50.261	NTAPIRet	NtSetInformationFile	cmd.exe	(0x8207E950) sr	SET_INFORMATION		SUCCESS, Inf
04:04:50.261	NTAPI	NtReadFile	cmd.exe	(0x8207E950) sr	READ	NORMAL	
04:04:50.261	Call	0x81F916D8-7		(0x8232E020) Ntfs	READ	NORMAL	
04:04:50.261	Call	0x821DF3C8-8		(0x8232E020) Ntfs	CLOSE		
04:04:50.261	Comp	0x821DF3C8-8		(0x8232E020) Ntfs	CLOSE		SUCCESS, Inf
04:04:50.261	Comp	0x81F916D8-7		(0x8232E020) Ntfs	READ		SUCCESS, Inf
04:04:50.261	NTAPIRet	NtReadFile	cmd.exe	(0x8207E950) sr	READ	NORMAL	SUCCESS, Inf
04:04:50.277	NTAPI	NtQueryInformationFile	cmd.exe	(0x8207E950) sr	QUERY_INFORMATION		
04:04:50.277	NTAPIRet	NtQueryInformationFile	cmd.exe	(0x8207E950) sr	QUERY_INFORMATION		SUCCESS, Inf
04:04:50.277	NTAPI	NtSetInformationFile	cmd.exe	(0x8207E950) sr	SET_INFORMATION		
04:04:50.277	NTAPIRet	NtSetInformationFile	cmd.exe	(0x8207E950) sr	SET_INFORMATION		SUCCESS, Inf
04:04:50.277	NTAPI	NtClose	cmd.exe	(0x8207E950) sr	CLEANUP		
04:04:50.277	Call	0x81F916D8-9		(0x8232E020) Ntfs	CLEANUP		
04:04:50.277	Comp	0x81F916D8-9		(0x8232E020) Ntfs	CLEANUP		SUCCESS, Inf
04:04:50.277	NTAPIRet	NtClose	cmd.exe	(0x8207E950) sr	CLEANUP		SUCCESS, Inf
04:04:50.277	NTAPI	NtQueryDirectoryFile	cmd.exe	(0x8207E950) sr	DIRECTORY_CONTROL	QUERY_DIRE...	
04:04:50.277	Call	0x81F916D8-10		(0x8232E020) Ntfs	DIRECTORY_CONTROL	QUERY_DIRE...	
04:04:50.277	Comp	0x81F916D8-10		(0x8232E020) Ntfs	DIRECTORY_CONTROL	QUERY_DIRE...	NO_MORE_FI
04:04:50.277	NTAPIRet	NtQueryDirectoryFile	cmd.exe	(0x8207E950) sr	DIRECTORY_CONTROL	QUERY_DIRE...	NO_MORE_FI
04:04:50.277	NTAPI	NtClose	cmd.exe	(0x8207E950) sr	CLEANUP		
04:04:50.277	Call	0x81F916D8-11		(0x8232E020) Ntfs	CLEANUP		

Windows Debugger Cheat Sheet

Command	Use
d	Dump memory address. Can further specify: a – ASCII d – Dword t - as type – provide a structure name to overlay (needs symbols) v - Variables local to scope (processes only) ps – Pointers and Symbols u – Unicode
k	Dump current thread stack
lm	List loaded and unloaded modules (useful to find drivers)
ln	List Nearest Symbols to address or symbol
r	Show/set registers
s	Search memory
u	Unassemble memory address or symbol. Also: ub (unassemble backwards from address/symbol)
.sympath+	Fix symbol path and append MS Symbol Server
.reload	Force reloading of Kernel Symbols

Useful Debugger Extensions:

Command	Use
!analyze	Crash dump analysis. The author's favorite ☺
!drvobj	Show Drive object of name (from lm)
!devobj	Show Device object at address
!devstack	Show Device Driver Stack for a given device
!idt	Show Kernel's Interrupt descriptor table (e.g. INT 2e)
!process	Show PEB at address (try "0")
!thread	Show TEB at address (try "0")
!pool, !pooltag, !poolfind	Pool debugging
!irpfind, !irp	Find IRPs in NonPagedPool, Display IRP contents

...If you liked this course, consider...

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Graceful introduction into the wonderful world of Linux for the non-command line oriented user. Basic skills and commands, work in shells, redirection, pipes, filters and scripting

Linux Administration:

Follow up to the Basic course, focusing on advanced subjects such as user administration, software management, network service control, performance monitoring and tuning.

Linux:

Linux User Mode Programming:

Programming POSIX and UNIX APIs in Linux, including processes, threads, IPC mechanisms and networking. Linux User experience required.

Linux Kernel Programming:

Guided tour of the Linux Kernel, 2.4 and 2.6, focusing on design, architecture, writing device drivers (character, block), performance and network devices

Embedded Linux Kernel Programming:

Similar to the Linux Kernel programming course, but with a strong emphasis on development on non-intel and/or tightly constrained embedded platforms

Windows Programming:

Windows Application Development, focusing on Processes, Threads, DLLs, Memory Management, and Winsock

Windows:

Windows Kernel Programming (this course):

Windows Kernel Architecture and Device Driver development – focusing on Network Device Drivers (in particular, NDIS) and the Windows Driver Model. Updated to include NDIS 6 and Winsock Kernel

Cryptography:

From Basics to implementations in 5 days: foundations, Symmetric Algorithms, Asymmetric Algorithms, Hashes, and protocols. Design, Logic and implementation

Security:

Application Security

Writing secure code – Dealing with Buffer Overflows, Code, SQL and command Injection, and other bugs... before they become vulnerabilities that hackers can exploit.